

Math 311

Numerical Methods

2.4: Error Analysis for Iterative Methods

Solutions of Equations of One Variable

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1 Introduction

- Suppose we have a sequence of real numbers $\{p_n\}_{n=0}^{\infty}$ that converge to a number p .
- In other words, $p_n \rightarrow p$.
- We want to talk about how **fast** the values are converging to p as n increases.
- Allowing us to talk about this enables us to make improvements and progress!
- We can then develop new methods and compare them to the older methods.
- We will be discussing the limit in the box below. Note that the top and bottom are how close the iterations are to p . Note also the bottom has a power of a positive α .

$$\lim_{n \rightarrow \infty} \frac{|p_{n+1} - p|}{|p_n - p|^\alpha} = ? \quad (1)$$

- The answer to this limit will tell us the speed of a sequence.
- What values are possible? (Discuss!)
 - ∞ (bad news).
 - 0 (good news)
 - somewhere in between $0 < \lambda < \infty$ (finite and positive) (best news)

Convergence a sequence of order α .

Definition. Suppose $\{p_n\}_{n=0}^{\infty}$ is a sequence that converges to p .

If constants $\lambda > 0$ and $\alpha > 0$ exist with

$$0 < \lim_{n \rightarrow \infty} \frac{|p_{n+1} - p|}{|p_n - p|^\alpha} = \lambda < \infty, \quad \text{then} \quad (2)$$

$\{p_n\}_{n=0}^{\infty}$ converges to p of order α , with asymptotic error constant λ .

Convergence examples

Definition. Further vocabulary on speeds of convergence include:

- If $\alpha = 1$ and $\lambda > 0$ then the sequence is said to converge linearly to p .
- If $\alpha = 2$ and $\lambda > 0$ then the sequence is said to converge quadratically to p .
- If $\alpha = 3$ and $\lambda > 0$ then the sequence is said to converge cubically to p .
- Similar language is used for larger values of α

Corollary. If $0 < \alpha < 1$ and $\lambda > 0$, the sequence converges sub-linearly to p .

Super-convergence of a sequence of order α .

Definition. Suppose $\{\mathbf{p}_n\}_{n=0}^{\infty}$ is a sequence that converges to \mathbf{p} .

If a positive constant $\alpha > 0$ exists with

$$\lim_{n \rightarrow \infty} \frac{|\mathbf{p}_{n+1} - \mathbf{p}|}{|\mathbf{p}_n - \mathbf{p}|^\alpha} = 0, \quad \text{then} \quad (3)$$

- $\{\mathbf{p}_n\}_{n=0}^{\infty}$ is said to converge to \mathbf{p} of order **super- α** .
- (or, equivalently \mathbf{p}_n **converges super- α -ly** to \mathbf{p}).
- (Note, the asymptotic error constant here **MUST** be $\lambda = 0$).

Super-convergence examples

Definition. Further vocabulary on speeds of convergence include:

- If $\alpha = 1$ and $\lambda = 0$ then \mathbf{p}_n is said to superlinearly converge to \mathbf{p} .
- If $\alpha = 2$ and $\lambda = 0$ then \mathbf{p}_n is said to super-quadratically converge to \mathbf{p} .
- If $\alpha = 3$ and $\lambda = 0$ then \mathbf{p}_n is said to super-cubically converge to \mathbf{p} .
- If $\alpha = 4$ and $\lambda = 0$ then \mathbf{p}_n is said to super-quartically converge to \mathbf{p} , etc.

Super- α -convergence implies existence of convergence of order β

Theorem. Suppose $\{p_n\}_{n=0}^{\infty}$ converges super- α -ly to p . It follows that:

- There exists $\beta > \alpha$ such that the asymptotic error constant, λ_β , is finite and positive. ($0 < \lambda_\beta < \infty$)
- In other words, **there exists a constant β larger than α such that p_n converges to p of order β**

Example.

- The Secant method is known for its superlinear (super-1-ly) convergence.
- This implies that it converges at a rate greater than 1.
- In fact, that rate is $\beta = \frac{1+\sqrt{5}}{2} \approx 1.618$.
- This means that once it gets close enough, it will improve upon the last iteration by increasing the number of digits correct by 61.8%.

Convergence of order α implies divergence for all $\beta > \alpha$

Theorem. Suppose that \mathbf{p}_n converges to \mathbf{p} of order α with $\lambda > 0$.

For every $\beta > \alpha$, \mathbf{p}_n does not converge to \mathbf{p} of order β . $\left(\lim_{n \rightarrow \infty} \frac{|\mathbf{p}_{n+1} - \mathbf{p}|}{|\mathbf{p}_n - \mathbf{p}|^\beta} = \infty \right)$

Proof. Let $\epsilon > 0$. Suppose $\beta = \alpha + \epsilon$. It follows that

$$\begin{aligned} \lim_{n \rightarrow \infty} \frac{|\mathbf{p}_{n+1} - \mathbf{p}|}{|\mathbf{p}_n - \mathbf{p}|^\beta} &= \lim_{n \rightarrow \infty} \frac{|\mathbf{p}_{n+1} - \mathbf{p}|}{|\mathbf{p}_n - \mathbf{p}|^{\alpha+\epsilon}} = \lim_{n \rightarrow \infty} \frac{|\mathbf{p}_{n+1} - \mathbf{p}|}{|\mathbf{p}_n - \mathbf{p}|^\alpha} \left(\frac{1}{|\mathbf{p}_n - \mathbf{p}|^\epsilon} \right) \\ &= \lim_{n \rightarrow \infty} \frac{|\mathbf{p}_{n+1} - \mathbf{p}|}{|\mathbf{p}_n - \mathbf{p}|^\alpha} \lim_{n \rightarrow \infty} \left(\frac{1}{|\mathbf{p}_n - \mathbf{p}|^\epsilon} \right) \\ &= \lambda \cdot \lim_{n \rightarrow \infty} \left(\frac{1}{|\mathbf{p}_n - \mathbf{p}|^\epsilon} \right) = \infty \end{aligned}$$

Therefore, it follows that $\{\mathbf{p}_n\}_{n=0}^\infty$ does not converge to \mathbf{p} of order $\beta > \alpha$. □

Convergence of order α implies superconvergence for all $\beta < \alpha$

Theorem. Suppose $\{\mathbf{p}_n\}_{n=0}^{\infty}$ converges to \mathbf{p} of order α . It follows that:

For every $\beta < \alpha$, \mathbf{p}_n converges to \mathbf{p} of order super- β .

Proof. Let $\epsilon > 0$. Suppose $\beta = \alpha - \epsilon$. It follows that

$$\begin{aligned}\lim_{n \rightarrow \infty} \frac{|\mathbf{p}_{n+1} - \mathbf{p}|}{|\mathbf{p}_n - \mathbf{p}|^\beta} &= \lim_{n \rightarrow \infty} \frac{|\mathbf{p}_{n+1} - \mathbf{p}|}{|\mathbf{p}_n - \mathbf{p}|^{\alpha - \epsilon}} = \lim_{n \rightarrow \infty} \frac{|\mathbf{p}_{n+1} - \mathbf{p}|}{|\mathbf{p}_n - \mathbf{p}|^\alpha} (|\mathbf{p}_n - \mathbf{p}|^\epsilon) \\ &= \lim_{n \rightarrow \infty} \frac{|\mathbf{p}_{n+1} - \mathbf{p}|}{|\mathbf{p}_n - \mathbf{p}|^\alpha} \lim_{n \rightarrow \infty} (|\mathbf{p}_n - \mathbf{p}|^\epsilon) \\ &= \lambda \cdot \lim_{n \rightarrow \infty} (|\mathbf{p}_n - \mathbf{p}|^\epsilon) = 0\end{aligned}$$

Therefore, it follows that \mathbf{p}_n converges to \mathbf{p} of order super- β , where $\beta = \alpha - \epsilon$. □

- So, in summary, for every sequence that converges of order α , we have one of the following three options:
 - For α , the sequence \mathbf{p}_n converges to \mathbf{p} of order α .
 - For all values less than α , the sequence converges super- α -ly to \mathbf{p} .
 - For all values greater than α , the sequence does not converge to \mathbf{p} .

1.1 Example 1

1. Let $p_n = \frac{1}{n^k}$ for some fixed $k > 0$. This sequence converges to $p = 0$. By the definition in (2),

$$\lambda = \lim_{n \rightarrow \infty} \frac{|p_{n+1} - p|}{|p_n - p|^\alpha} = \lim_{n \rightarrow \infty} \frac{1}{\left[\frac{1}{n^k}\right]^\alpha} = \lim_{n \rightarrow \infty} \left[\frac{n^\alpha}{n+1}\right]^k = \left[\lim_{n \rightarrow \infty} \frac{n^{\alpha-1}}{1 + \frac{1}{n}}\right]^k = \left[\lim_{n \rightarrow \infty} n^{\alpha-1}\right]^k$$

It follows that

- If $\alpha < 1$ then $\lambda = 0$, which implies super- α convergence.
- If $\alpha = 1$ then $\lambda = 1$, which implies the sequence converges **linearly**.
- If $\alpha > 1$ then $\lambda = \infty$, which implies it does **not** converge at any rate greater than linear.

1.2 Example 2

2. Let $p_n = 10^{-2^n}$. This also converges to 0 and the asymptotic error constant is

$$\lambda = \lim_{n \rightarrow \infty} \frac{|p_{n+1} - p|}{|p_n - p|^\alpha} = \lim_{n \rightarrow \infty} \frac{10^{-2^{n+1}}}{10^{-\alpha \cdot 2^n}} = \lim_{n \rightarrow \infty} 10^{-2 \cdot 2^n} \cdot 10^{\alpha \cdot 2^n} = \lim_{n \rightarrow \infty} 10^{(\alpha-2) \cdot 2^n}$$

- If $\alpha < 2$, then $\lambda = 0$. This means that p_n converges super-linearly to p .
- If $\alpha = 2$, then $\lambda = 1$. This means that p_n converges quadratically to p .
- If $\alpha > 2$, then $\lambda = \infty$ and p_n does not converge at a rate larger than $\alpha = 2$.

Using this as a typical quadratic sequence, you can see how fast a quadratically convergent sequence moves. The terms of the sequence are:

n	p_n
1	10^{-2}
2	10^{-4}
3	10^{-8}
4	10^{-16}
5	10^{-32}
6	10^{-64}

It doubles the number of correct digits with each iteration! Note that after 6 iterations, it is already really small!

2 Fixed Point Iteration (Section 2.2)

How good is fixed point iteration? Let's analyze the fixed point algorithm,

$$p_{n+1} = g(p_n)$$

with fixed point p . The key to the speed of convergence is derivatives of $g(p)$.

Convergence of Fixed Point Iteration:

Theorem. Let $g \in C[a, b]$ and $g' \in C(a, b)$. Furthermore,

assume there exists $k < 1$ such that $|g'(x)| \leq k$ for all x in (a, b) .

- If $g'(p) \neq 0$, the sequence converges linearly to the fixed point p .
- If $g'(p) = 0$, the sequence converges at least quadratically to the fixed point p .

Proof. • First, we will show that $p_n \rightarrow p$. Start with the statement

$$p_n = g(p_{n-1})$$

- Subtract p from both sides and take the absolute value:

$$|p_n - p| = |g(p_{n-1}) - p|$$

- Note that since $g(\mathbf{p}) = p$, it follows that

$$|\mathbf{p}_n - \mathbf{p}| = |g(p_{n-1}) - g(\mathbf{p})|$$

- Suppose that $|g'(x)| \leq k < 1$. By the MVT, ξ exists between p_{n-1} and \mathbf{p} where

$$\underbrace{|\mathbf{p}_n - \mathbf{p}|}_{\text{left side}} = |g(p_{n-1}) - g(\mathbf{p})| = |g'(\xi)||p_{n-1} - p| \leq k \underbrace{|p_{n-1} - p|}_{\text{right side}}$$

- Thus, the left side is less than or equal to the right side!

$$|\mathbf{p}_n - \mathbf{p}| \leq k|p_{n-1} - p|. \quad (4)$$

- Recursively plugging equation (4) into itself yields:

$$|\mathbf{p}_n - \mathbf{p}| \leq k|k|p_{n-2} - \mathbf{p}|| \leq k|k|k|p_{n-3} - \mathbf{p}|| \leq \cdots \leq k^n|p_0 - p|.$$

- Since $0 \leq k < 1$, then $k^n \rightarrow 0$ as $n \rightarrow \infty$.
- Therefore, $|\mathbf{p}_n - \mathbf{p}| \rightarrow 0$ as $n \rightarrow \infty$ (which is equivalent to $\mathbf{p}_n \rightarrow p$).
- Next, it can be shown that convergence speed will depend on the derivatives of g .
- By the Mean Value Theorem, ξ exists between p_{n-1} and \mathbf{p} where

$$|\mathbf{p}_{n+1} - \mathbf{p}| = |g(p_n) - g(\mathbf{p})| = |g'(\xi_n)||p_n - p|$$

- Since $p_n \rightarrow p$, then $\xi_n \rightarrow p$ as $n \rightarrow \infty$. Thus,

$$\lim_{n \rightarrow \infty} \frac{|p_{n+1} - p|}{|p_n - p|} = \lim_{n \rightarrow \infty} |g'(\xi_n)| = \left| g' \left(\lim_{n \rightarrow \infty} \xi_n \right) \right| = |g'(p)|$$

- It follows that

$$\lim_{n \rightarrow \infty} \frac{|p_{n+1} - p|}{|p_n - p|} = |g'(p)| \quad (5)$$

- By the definition of convergence and super-convergence in (2) and (3), we see that:
 - ⎧ If $|g'(p)| \neq 0$, then p_n converges to p at a linear rate ($\lambda = |g'(p)| > 0$).
 - ⎧ If $|g'(p)| = 0$, then p_n converges to p at a **super-linear** rate.
- To precisely find the order that p_n converges to p , expand $g(x)$ in a Taylor's Polynomial about p . It follows that by Taylor's Theorem,

$$g(x) = g(p) + g'(p)(x - p) + \frac{1}{2}g''(\xi)(x - p)^2, \quad (6)$$

where ξ is between x and p .

- Since $g(\mathbf{p}) = p$ and $g'(\mathbf{p}) = 0$, then it follows that (6) simplifies as

$$g(x) = g(\mathbf{p}) + \cancel{g'(\mathbf{p})}^0(x - \mathbf{p}) + \frac{1}{2}g''(\xi)(x - \mathbf{p})^2$$

$$g(x) = p + \frac{1}{2}g''(\xi)(x - \mathbf{p})^2 \quad (7)$$

- Evaluate (7) at $x = \mathbf{p}_n$. Note when $x = \mathbf{p}_n$, $g(\mathbf{p}_n) = \mathbf{p}_{n+1}$,

$$g(\mathbf{p}_n) = p + \frac{1}{2}g''(\xi_n)(\mathbf{p}_n - \mathbf{p})^2$$

$$g(\mathbf{p}_n) - p = \frac{1}{2}g''(\xi_n)(\mathbf{p}_n - \mathbf{p})^2$$

$$\frac{|\mathbf{p}_{n+1} - p|}{|\mathbf{p}_n - \mathbf{p}|^2} = \frac{1}{2}|g''(\xi_n)|$$

- \mathbf{p}_n converges to \mathbf{p} . Further, since ξ_n is always between \mathbf{p}_n and \mathbf{p} , then ξ_n converges to \mathbf{p} as well. It follows that

$$\lim_{n \rightarrow \infty} \frac{|\mathbf{p}_{n+1} - p|}{|\mathbf{p}_n - \mathbf{p}|^2} = \frac{1}{2} \lim_{n \rightarrow \infty} |g''(\xi_n)| = \frac{1}{2} \left| g'' \left(\lim_{n \rightarrow \infty} \xi_n \right) \right| = \frac{1}{2} |g''(\mathbf{p})|$$

- This yields a similar statement to (5).

$$\lim_{n \rightarrow \infty} \frac{|p_{n+1} - p|}{|p_n - p|^2} = \frac{1}{2}|g''(p)|$$

- We then can conclude that

{ If $|g''(p)| > 0$, then p_n converges to p of order 2 (quadratically).
{ If $g''(p) = 0$, then p_n converges to p at a super-quadratic rate (at least quadratic).

□

2.1 Final Thoughts

- A side note: a key to finding faster fixed point methods is to generate a g for which $g'(p) = 0$. This leads to a quadratic rate.
- We can continue this idea and develop faster methods by finding a function $g(x)$ for which $g'(p) = g''(p) = 0$, but $g'''(p) \neq 0$. This will lead to cubic convergence rate.
- There are some examples of these methods, one called “Halley’s Method” and another “Olvert’s Method”. We get a faster speed by in exchange with more complexity. There is also a series of methods called “Householder’s Methods” which can generate sequences of any desired rate.

3 Newton's Method

- We will develop a faster fixed point method using the tricks above.
- We want to find $g(x)$ such that $g'(p) = 0$, where p is the fixed point.
- In the past, to solve $f(x) = 0$, we created $g(x) = x - f(x)$, or something similar.
- Let's assume that the form of $g(x)$ is as follows and that we want to find $\phi(x)$ which forces $g'(p) = 0$.

$$\begin{aligned}g(x) &= x - \phi(x)f(x) \\g'(x) &= 1 - \phi'(x)f(x) - \phi(x)f'(x) \quad \text{(Using the product rule)}\end{aligned}$$

- Since $f(p) = 0$ and forcing $g'(p) = 0$, it follows that

$$0 = g'(p) = 1 - \phi'(p)f(p) - \phi(p)f'(p) \quad \implies \quad \phi(p) = \frac{1}{f'(p)}$$

- Therefore, the function to iterate is:

$$\begin{aligned}g(x) &= x - \phi(x)f(x) = x - \left(\frac{1}{f'(p)}\right) f(x) \\g(x) &= x - \frac{f(x)}{f'(p)}\end{aligned}$$

- This is VERY similar to Newton's method, except we don't know $f'(\mathbf{p})$.
- Since \mathbf{p} is usually unknown, then let $\mathbf{p} = \mathbf{p}_n$ and then

$$\mathbf{p}_{n+1} = \mathbf{g}(\mathbf{p}_n) = \mathbf{p}_n - \frac{\mathbf{f}(\mathbf{p}_n)}{\mathbf{f}'(\mathbf{p}_n)},$$

which IS Newton's Method.

- We can show quadratic convergence by analyzing the derivative of $g(x)$ at \mathbf{p} .
- Taking the derivative of this function $g(x)$ is

$$g'(x) = 1 - \frac{[f'(x)]^2 - f(x)f''(x)}{[f'(x)]^2} = \frac{f(x)f''(x)}{[f'(x)]^2} \quad (8)$$

- Since $f(\mathbf{p}) = 0$, then it is clear that $g'(\mathbf{p}) = 0$. Thus, Newton's method will converge at least quadratically!
- However, if $\mathbf{f}'(\mathbf{p}) = \mathbf{0}$ at the same time as $g'(\mathbf{p})$, then we might not have quadratic convergence.

3.1 Problems with Newton's Method

- Here's an example of the problem:

Let $f(x) = x^2$, so $f'(x) = 2x$. This has the obvious solution of 0. So $g(x)$ is

$$g(x) = x - \frac{x^2}{2x} \implies g(x) = \frac{x}{2}$$

- Note that $g'(x) = \frac{1}{2} \neq 0$
- All this method does it repeatedly half the answer from what was there before. This is clearly a linearly converging sequence. (Sounds like the Bisection)
- Why did this fail to have quadratic convergence?

Zero of Multiplicity m

Definition. A solution \mathbf{p} of $f(x) = 0$ is said to be a zero of multiplicity m of f if $f(x)$ can be written as

$$f(x) = (x - \mathbf{p})^m q(x), \text{ for } x \neq p, \text{ where } \lim_{x \rightarrow p} q(x) \neq 0.$$

Zero of Multiplicity m (Part 2)

Theorem.

The function $f \in C^m[a, b]$ has a zero of multiplicity m at \mathbf{p} if and only if

$$f(\mathbf{p}) = f'(\mathbf{p}) = f''(\mathbf{p}) = \dots = f^{(m-1)}(\mathbf{p}) = 0, \quad \text{but} \quad f^{(m)}(\mathbf{p}) \neq 0.$$

3.2 Example

The function $f(x) = 2 \cos x - 2 - x^2$ has a zero of multiplicity 2 at $x = 0$. Here is why:

$$\begin{aligned} f(x) &= 2 \cos x - 2 - x^2 & \implies & f(0) = 2(1) - 2 - 0^2 & = 0 \\ f'(x) &= -2 \sin x - 2x & \implies & f'(0) = -2(0) - 2(0) & = 0 \\ f''(x) &= -2 \cos x - 2 & \implies & f''(0) = -2(1) - 2 & = -4 \neq 0 \end{aligned}$$

- When the zero is NOT simple, then Newton's method will converge linearly.

4 Modified Newton's Method

- What can we do to fix this problem, if possible?
- There are two methods to fix this:
 - $g(x) = x - \frac{mf(x)}{f'(x)}$. (Problem 8 in Section 2.4, 5th Ed.)
 - Apply Newton's Method to $\mu(x) = \frac{f(x)}{f'(x)}$.
- The first method requires knowledge of the multiplicity of the root. This information is not available in general.
- The second method has the advantage of not requiring knowledge of m .
- Let's focus on the second method.

- Suppose $f(x) = (x - \mathbf{p})^m q(x)$, where $\lim_{x \rightarrow p} q(x) \neq 0$.
- It follows that $f'(x) = m(x - \mathbf{p})^{m-1} q(x) + (x - \mathbf{p})^m q'(x)$
- The new function $\mu(x)$ simplifies to

$$\begin{aligned} \mu(x) &= \frac{f(x)}{f'(x)} = \frac{(x - \mathbf{p})^m q(x)}{m(x - \mathbf{p})^{m-1} q(x) + (x - \mathbf{p})^m q'(x)} \\ &= \frac{(x - \mathbf{p})^m q(x)}{(x - \mathbf{p})^{m-1} [m q(x) + (x - \mathbf{p}) q'(x)]} && \text{(Factor out } (x - \mathbf{p})^{m-1} \text{)} \\ &= \frac{(x - \mathbf{p}) q(x)}{m q(x) + (x - \mathbf{p}) q'(x)} && \text{(Cancel factor of } (x - \mathbf{p})^{m-1} \text{)} \end{aligned}$$

- What remains is a simple root of $\mu(x)$ at $x = p$.
- So apply Newton's to $\mu(x) = \frac{f(x)}{f'(x)}$!
- We will now ignore the form of $f(x)$ above, and keep it general.

- So it follows

$$\begin{aligned}
 g(x) &= x - \frac{\mu(x)}{\mu'(x)} = x - \frac{\frac{f(x)}{f'(x)}}{\frac{d}{dx} \left(\frac{f(x)}{f'(x)} \right)} \\
 &= x - \frac{\frac{f(x)}{f'(x)}}{\frac{f'(x) \cdot f'(x) - f(x) f''(x)}{[f'(x)]^2}} \\
 &= x - \frac{f(x)}{f'(x)} \left(\frac{[f'(x)]^2}{f'(x) \cdot f'(x) - f(x) f''(x)} \right) \\
 g(x) &= x - \frac{f(x) f'(x)}{[f'(x)]^2 - [f(x)] [f''(x)]} \tag{9}
 \end{aligned}$$

- It follows that the “Modified Newton’s Method” is

$$\mathbf{p}_{n+1} = \mathbf{p}_n - \frac{f(\mathbf{p}_n) f'(\mathbf{p}_n)}{[f'(\mathbf{p}_n)]^2 - [f(\mathbf{p}_n)] [f''(\mathbf{p}_n)]} \tag{10}$$

4.1 Example 1 of Modified Newton’s Method

- The previous example of a slow down was applying Newton’s Method to $f(x) = x^2$, which lead to $g(x) = \frac{1}{2}x$.

- Let's try the modified procedure on this function now.
- So $f'(x) = 2x$ and $f''(x) = 2$. The Modified Newton's simplifies to

$$\begin{aligned}g(x) &= x - \frac{f(x)f'(x)}{[f'(x)]^2 - [f(x)][f''(x)]} = x - \frac{x^2(2x)}{[2x]^2 - [x^2][2]} \\ &= x - \frac{2x^3}{4x^2 - 2x^2} = x - \frac{2x^3}{2x^2} = x - x = 0\end{aligned}$$

- Woah! That is really sped up! The guess is always 0! It found the root in 1 step!
- Any function like $f(x) = (x - c)^2$ will converge in 1 iteration to $x = c$ using the modified Newton's Method.

4.2 Example 2 of Modified Newton's Method

- Apply it to a similar problem: $f(x) = x^3 - 3x + 2$ has a double root at $x = 1$
- Newton's Method applied to it is $g(x) = x - \frac{x^3 - 3x + 2}{3x^2 - 3}$.
- With an initial guess of $x_0 = 2$, it takes 35 iterations to converge within 10^{-13} .
- The Modified Newton's simplifies to $(f'(x) = 3x^2 - 3$ and $f''(x) = 6x)$

$$g(x) = x - \frac{f(x)f'(x)}{[f'(x)]^2 - [f(x)][f''(x)]} = x - \frac{(x^3 - 3x + 2)(3x^2 - 3)}{[3x^2 - 3]^2 - [x^3 - 3x + 2][6x]} = \frac{4x + 2}{x^2 + 2x + 3}$$

- With an initial guess of $x_0 = 2$, it takes 4 iterations to converge within 10^{-10} .

4.3 Example with $f(x) = 2 \cos x - 2 - x^2$

- In example 2.3, we used $f(x) = 2 \cos x - 2 - x^2$.
- It had a zero of multiplicity 2 at $x = 0$.
- Starting with a guess of $x_0 = 1$, we applied Newton's and Modified Newton's to it.
- It takes 37 iterations to converge within 10^{-8} with Newton's.
- It takes 7 iterations to converge within 10^{-8} with Modified Newton's.

4.4 Orders of Common Methods

Method	Iteration Formula ($\mathbf{p}_{n+1} =$) or Combination	Worst Order	Best Order	Global Convergence?
Bisection	Not iteration	1	1	Yes
Fixed Point	$\mathbf{p}_{n+1} = g(\mathbf{p}_n)$	ϵ	–	No
Newton's	$\mathbf{p}_{n+1} = g(\mathbf{p}_n) = \mathbf{p}_n - \frac{f(\mathbf{p}_n)}{f'(\mathbf{p}_n)}$	1	2	No
Steffensen's	Fixed Pt & Aitkens	2	2	No
Modified Newton's I	$\mathbf{p}_{n+1} = g(\mathbf{p}_n) = \mathbf{p}_n - \frac{f(\mathbf{p}_n)f'(\mathbf{p}_n)}{[f'(\mathbf{p}_n)]^2 - f(\mathbf{p}_n)f''(\mathbf{p}_n)}$	2	2	No
Modified Newton's II	$g(\mathbf{p}_n) = \mathbf{p}_n - m \frac{f(\mathbf{p}_n)}{f'(\mathbf{p}_n)}$	2	2	No
Secant	$\mathbf{p}_{n+1} = g(\mathbf{p}_n) = p_n - \frac{f(\mathbf{p}_n)(\mathbf{p}_n - p_{n-1})}{f(\mathbf{p}_n) - f(p_{n-1})}$	–	1.618	No
False Position	hybrid	1	1.618	Yes
Illinois	hybrid	–	1.442	Yes
Halley's Method	$\mathbf{p}_{n+1} = g(\mathbf{p}_n) = \mathbf{p}_n - \frac{f(\mathbf{p}_n)f'(\mathbf{p}_n)}{[f'(\mathbf{p}_n)]^2 - \frac{1}{2}f(\mathbf{p}_n)f''(\mathbf{p}_n)}$	–	3	No

4.5 Other Methods we didn't cover

Method	Iteration Formula ($p_{n+1} =$) or Combination	Worst Order	Best Order	Global Convergence?
Brent's	hybrid	1.618	1.839	Yes
IQI	Inverse Quadratic Interpolation		1.839	No
ITP Method	Interpolate, Truncate, and Project	1	> 1	Yes
Mueller's Method	secant & IQI		1.839	?
Laguerre's Method	general poly root solver	1	3	Almost
Jenkins-Traub Method	complete polynomial root solver	1	2.618	Yes
Ridder's Method	false position variant	1.414	2	Yes
Durand-Kerner Method	simultaneously all roots of polynomial	1	2	Yes
Aberth Method	simultaneously all roots of polynomial	1	3	Yes